SAP: Stall-Aware Pacing for Improved DASH Video Experience in Cellular Networks

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ABSTRACT

The dramatic growth of cellular video traffic represents a practical challenge for cellular network operators in providing a consistent streaming Quality of Experience (QoE) to their users. Satisfying this objective has so-far proved elusive, due to the inherent system complexities that degrade streaming performance, such as variability in both video bitrate and network conditions. In this paper, we present SAP as a DASH video traffic management solution that reduces playback stalls and seeks to maintain a consistent QoE for cellular users, even those with diverse channel conditions. SAP achieves this by leveraging both network and client state information to optimize the pacing of individual video flows. We extensively evaluate SAP performance using real video content and clients, operating over a simulated LTE network. We implement state-of-the-art client adaptation and traffic management strategies for direct comparison. Our results, using a heavily loaded base station, show that SAP reduces the number of stalls and the average stall duration per session by up to 95%. Additionally, SAP ensures that clients with good channel conditions do not dominate available wireless resources, evidenced by a reduction of up to 40% in the standard deviation of the QoE metric.

CCS CONCEPTS

• **Information systems** → **Multimedia streaming**; • **Networks** → *Network performance modeling*; *Network simulations*;

KEYWORDS

Adaptive bitrate video streaming, DASH, QoE, separable programming

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1 INTRODUCTION

Video streaming over cellular networks is growing rapidly, and is expected to reach 75% of all cellular traffic by 2020[1]. This rapid growth has network operators playing catch-up trying to ensure that all of their customers are given a consistent Quality of Experience (QoE). In cellular networks, offering quality assurances is especially challenging due to the complexities of the wireless transmission medium. In a recent study by Conviva[2], the majority of users identified video stalls as the most irritating factor while streaming. Hence, it is surprising that approximately 50% of video sessions encounter stalls[3]. There is a crucial need to develop a deep understanding of the cause of video stalls, and to design techniques to minimize stalls.

Video streaming today involves a client player fetching segments from a server over HTTP. The most popular standard is Dynamic Adaptive Streaming over HTTP (DASH). The client continuously adapts the selected video quality based on prevailing network conditions. A well-known issue with DASH is when multiple video users compete in sharing a network link, with several studies [1, 13, 15] demonstrating that quality instability and stalls are very common. Such issues are greatly exacerbated in wireless networks. To illustrate, when multiple users share a cellular base station, there exist at least three interacting control loops involved in the streaming process. The DASH client exercises the control on the video quality control loop with the server, the streaming server and

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client are involved in the underlying TCP congestion control loop, and the base station scheduler controls the resource allocation and scheduling over the downlink air interface control loop to the cellular device. Each of these control loops operates independently and at different timescales, which often leads to degraded streaming performance [4, 7].

Techniques proposed to improve streaming performance consider different approaches including end-to-end [2, 5, 6, 11, 15, 16], network-based [4, 7, 13, 18, 22, 25], and hybrid solutions [3, 10, 18, 19]. In end-to-end solutions, changing the client adaptation logic and/or the server delivery behavior represent common design elements. In network-based solutions, traffic shaping, trans-rating, and/or trans-coding are examples of the functions considered. These approaches maintain the independence between network operators and content providers, but they may lead to sub-optimal performance. For example, a streaming client may trade video quality by adopting an overly conservative estimator for the available network throughput to reduce stalls. Similarly, network-based solutions usually consider network-state only in their decisions and the application state is overlooked. Hence, a video client may still suffer from performance degradation, such as video stalls. Hybrid solutions assume interaction between end-nodes and network agents. Additionally, these solutions may assume integrated control loops for both the end-to-end level and the network level by operating with a bird's eye-view. But this comes at the cost of communication overhead. The MPEG Server and Network Assisted DASH (SAND) standard provides such a framework, with a DASH-Aware Network Element (DANE) as a network agent.

In delivering video over a wired link, the goal is to allocate resources in a fair manner so that clients can have equivalent QoE. But the nature of wireless access means that the achievable data rate is not the same for each client, as it is a function of the channel condition at each client. In addition, each client's channel condition can vary considerably over short time periods due to fading and other factors. Thus, unlike in wired links, it must be recognized that achieving equivalent QoE across all clients is an unreasonable objective. Fairness is not absolute, but is rather a function of the channel conditions available to each client. Thus, steps must be taken to manage wireless resources to ensure that clients with poorer channel conditions do not suffer unnecessarily, while clients with good channel conditions always seek to maximize their own QoE. This can be achieved by redistributing channel resources so as to reduce extremes of QoE across clients. Furthermore, the potential of a stall impairs user QoE significantly, and it is desirable for the network to re-allocate resources to such a client on a short term basis, without introducing substantial unfairness. Our approach is to guide this redistribution process while balancing multiple considerations, including that of avoiding stalls.

In this research, we set out to see if we can eliminate stalls (or come close to it) for cellular users through the use of a network agent that is aware of a limited amount of client state, and can judiciously manage resources at the bottleneck wireless link. We further seek to design a solution that can manage the link resources such that the Quality of Experience across all clients is well-balanced, thus promoting fairness. We introduce Stall Aware Pacing (SAP) as a novel network-based solution to improve the streaming performance in cellular systems. SAP indirectly integrates the client-server quality control loop and the base station downlink scheduling process to reduce stalls when a group of users share a congested cellular downlink air interface. SAP achieves this goal by optimizing the delivery rate of individual packets of a flow based on both application and network states, targeting resources to clients for whom a stall is imminent. These clients are often, but not always, those at the cell edge. Our performance evaluation based on real video sessions (H.264 video streamed from a server to clients playing the video) over a simulated LTE network shows that SAP improves the stall performance for different DASH client adaptation algorithms in a wide variety of operating conditions in comparison to state-of-the-art techniques. Moreover, SAP achieves this with only a modest reduction in the average delivered video rate. Our contributions can be summarized as follows:

- We developed SAP as a novel network-based optimized pacing solution whose design captures the typical user perception for image quality while factoring in the impact of stalls. Additionally, the SAP design provides for differential user treatment to accommodate inherent design features such as device capability or user priority.
- We present collaborative SAP, in which the state of the client buffer is relayed by the client to the network-based pacing module. We also present non-collaborative SAP, an in-network algorithm that operates on an estimate of the buffer-level at the client.
- We evaluate SAP using a laboratory test-bed with video clients streaming real video content over a simulated LTE network in a wide variety of scenarios, including different user topologies and fading conditions.
- We show that SAP reduces the number of stalls and the average stall duration by up to 95% leading to a dramatic reduction in the stall QoE penalty, up to 85%. We show that SAP reduces the variability in the client QoE by up to 40%, being especially effective when clients operate with different network conditions. Finally, we conclude that the extra requirement for access to client state information in collaborative SAP does not provide a noticeable benefit over non-collaborative SAP.

The rest of this paper is organized as follows. Background and related work is presented in Section 2 followed by the design of SAP in Section 3. We then present our performance evaluation setup and results in Section 4. Conclusions and future work are presented in Section 5.

2 BACKGROUND AND RELATED WORK

Adaptive streaming over HTTP, recently standardized as DASH, is becoming the dominant technique for transmitting

video due to its ability to traverse firewalls and the abundance of HTTP infrastructure. With DASH, the video is split into multiple segments and each segment is encoded into different representations varying in their qualities. DASH video clients may change the video quality at segment boundaries in response to variations in operating conditions. Client adaptation strategies span different approaches including buffer-based, rate-based, and hybrid approaches.

Huang et al. [14] propose a buffer-based strategy by which the buffer-level is mapped to the selected video quality. Jiang et al. [15] propose FESTIVE as a rate based heuristic with randomized scheduling of segment requests and stateful adaptive rate update strategy. In [5], De Cicco et al. propose ELASTIC in which segment quality selection is based on a proportional integral controller. In [16], Li et al. propose PANDA that employs self-traffic network probing to establish an accurate estimate of the available network throughput. These adaptation strategies include design elements to counter stalls such as maintaining a high level of playout buffer occupancy and/or conservative rate estimators. However, achieving the best streaming performance in the highly variable operating conditions of a cellular network is difficult to achieve while solely relying on client adaptation.

A few server-based techniques have been proposed to improve the streaming performance. In [6], De Cicco et al. propose a closed-loop controller implemented at both the client and server to take adaptation decisions and regulate packet transmissions from the server. In [11], Ghobadi et al. propose Trickle as a server-side solution that enforces an upper-bound on the TCP congestion window. This bound is adjusted according to the streaming rate and round-trip time. However, such solutions are agnostic to the underlying operating conditions and are not designed to achieve the best performance in highly variable bandwidth conditions observed in cellular networks, which is the focus of our work.

A number of network-based solutions proposed to improve network resource sharing among multiple video clients are more directly related to our work. In [13], Houdaille and Gouache show that video rate shaping at a WiFi home gateway reduces the number of quality changes and oscillations for both Microsoft smooth streaming and Apple's HLS. In [22], Pu et al. employ a proxy at the edge of network core. This proxy implements split TCP, rate-dependent packet prioritization, and video transrating based on optimizing a system objective function integrating user rate-utility, smooth rate switching, and buffer-level tracking. [28] takes advantage of mobile-CDN, as part of the network operator infrastructure, to implement an application-level fair scheduler for video rate control. The proposed scheduler integrates user, device, network, and TCP information in its decisions. In [7], a target rate for every data and video stream is communicated to an underlying minimum rate proportional scheduler at the base station. This target rate is estimated using an adaptive guaranteed bit rate algorithm that is shown to match the optimal solution of maximizing the total user utility under resource constraints. In [4], Chen et al. propose the AVIS scheduling framework that throttles each stream to a rate in a specific range. The minimum rate in the range is identified by solving an optimization problem that maximizes the total user utility in the wireless system assuming limited resources. The maximum rate is determined based on the allocated minimum rate and the next higher video rate. A crucial point of distinction for our SAP solution is that we adopt a QoE-driven utility that integrates stall probability in its decision, and the ability to incorporate device heterogeneity as a factor. Furthermore, the aforementioned network-based solutions do not allow one to benefit from state updates from clients, thus distinguishing SAP, as it can operate whether or not such collaboration is available.

Collaborative solutions assume some level of interaction between clients and network elements. In [18], Mok et al. propose using a network proxy in the content provider network to assist the client in measuring the available bandwidth by monitoring packet round trip times. The client integrates this estimate in its QoE-driven adaptation policy. In [25], the authors show that maximizing the minimum buffer-level for multiple non-adaptive video clients in a wireless system is an NP-hard problem. With the client providing the buffer-level at every epoch, they propose a greedy scheduler for non-adaptive video and show its optimality if the wireless medium remains stable between decision epochs. In [3], Bouten et al. propose an in-network QoE-driven quality selection based on both network and client information. The wired network resources are monitored using a packet sampling approach to accurately forecast future available resources in wired networks. The network then solves an optimization problem that maximizes a total video QoE metric that includes received quality, the number of rate switches and stalls, subject to resource constraints. The result is passed directly to the client which has been altered to use it in selecting the quality of the next segment. Georgopoulos et al. [10] propose a framework to identify individual user quality for fair resource allocation among a group of users. The estimated rate is then communicated directly back to the customized client using the northbound interface of the software defined network controller. In contrast, SAP uses the solution of its optimization to *indirectly* impact the client behavior, not requiring that the clients be altered to accept directions from the network. Thus, SAP also maintains the independence of control functions of both network and content providers, which from a practical viewpoint is advantageous for companies operating in this highly competitive business sector.

3 SAP

3.1 SAP Overview

SAP is a stall-aware network element that manages a group of video flows and seeks to minimize stalls, while improving the QoE of video clients sharing a cellular base station (BS). SAP achieves this goal by pacing the delivery rate of video packets while considering both client application and network state. Network state is captured by the amount of resources dedicated to clients receiving video traffic, the number of

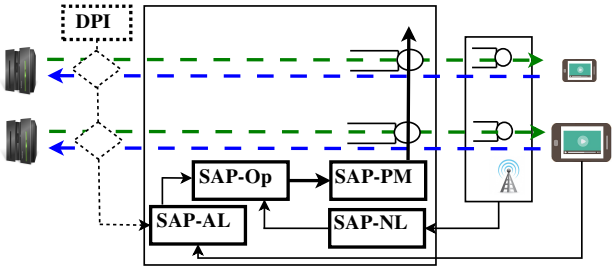


Figure 1: SAP Architecture, comprising the application state logger (SAP-AL), network state logger (SAP-NL), SAP pacing optimizer (SAP-Op), and SAP pacing manager (SAP-PM)

clients and the channel quality information for individual clients. Application state information includes video encoding information and the current video buffer-level at the client. Figure 1 shows the SAP components.

SAP may operate in either a collaborative or non-collaborative mode. In collaborative mode, the streaming client provides the application logger the current application state information. In the non-collaborative mode, SAP has to adopt different techniques (e.g., deep packet inspection) to obtain or estimate the application state information. SAP's network state logger interacts with network entities to collect relevant network state information including the number of DASH clients, their corresponding channel quality information, and the amount of resources dedicated for them at the BS. The pacing optimizer integrates the application and network state information to determine the best delivery rate for each individual video flow. SAP's pacing manager controls the delivery rate for each flow based on the rate calculated by the optimizer.

3.2 Design of the SAP Pacing Optimizer

Typically, the adaptation algorithm of DASH determines the quality of the next segment it is going to request. The time between consecutive requests is therefore of the order of a segment duration, which is typically between 4-10 seconds for cellular networks. SAP's optimization may execute periodically, or it may be executed in response to a predefined set of triggers such as client departure, client arrival, change in network resources and/or similar related events. SAP's optimizer could also be invoked in a hybrid combination of both periodic and explicit trigger events. The periodic execution would be performed at a much longer time scale than the BS scheduling time. The cellular BS scheduler executes every few milliseconds, and in LTE it allocates resources every two milliseconds. The SAP optimizer would run at a frequency in between this BS scheduler frequency and the client application adaptation frequency of once every 4-10 seconds. We expect a typical period between 250 ms and 1 second as being suitable for performing this pacing optimization. Such period would enable SAP to perform several rate adjustments per segment to orchestrate the interaction between different streams without incurring excessive unnecessary processing overhead.

3.2.1 System Model. We consider a base station (BS) serving U DASH users that each have a different channel quality identified by its spectral efficiency per resource unit (RU) $\gamma_u kbps/RU$, where $u \in \{1..U\}$. We assume that the BS capacity is divided into a group of allocatable resource units managed by the BS scheduler. In LTE, a resource block group (RBG) at an eNodeB represents the allocatable resource unit. The size of a RBG would vary depending on the BS bandwidth. For example, the RBG is a single resource block for a BS with 1.4MHz cell bandwidth while 2-resource blocks form a single RBG in a BS with bandwidth of 3MHz or 5Mhz. As the BS bandwidth increases, RBG are formed from larger resource block clusters to simplify the scheduler design. We assume that C RUs are dedicated for DASH users. Generally, C would vary dynamically depending on different factors, including the volume of non-video traffic that may be sharing the BS with DASH clients. We assume the presence of a bandwidth slicer that distributes the cell bandwidth among different traffic classes [4], e.g., DASH video, non-DASH video, and background traffic. In compliance with DASH, we consider each video is split into S segments with each segment corresponding to a duration of τ seconds. The videos are encoded into Q quality representations whose average rates are denoted r_q, where $q \in \{1..Q\}$.

3.2.2 SAP QoE-Oriented Design. The SAP pacing optimizer integrates stalls and video quality as key aspects affecting the QoE performance of video delivery over cellular systems. Specifically, SAP maximizes a video quality utility metric and minimizes a stall penalty that is a significant factor for user QoE. SAP's rate changes are presented to the client adaptation logic in a manner similar to changes in the channel condition. Our evaluation confirms that SAP does not adversely impact the end-client switching performance and in many cases reduces the number of quality switches.

SAP captures the visual quality using a tunable concave quality utility metric inspired by the video quality metric (VQM) [21]. For the same content encoded at different rates, VQM varies between 0 to 1, with higher quality having a lower value. Further, it is well understood that as the encoding rate increases, the marginal improvement in quality reduces. Therefore, we consider an exponential video quality utility measure, denoted by $\Upsilon_u(x_u)$, expressed as

$$\Upsilon_u(x_u) = (1 - e^{-\rho_u x_u / r_q}), \quad (1)$$

where x_u is the rate that would be allocated to user u and ρ_u is a tunable parameter that can be set according to the device capability or based on operator requirements. Note that in Eq. (1), higher rates have larger utility. Also, $\Upsilon_u(x_u)$ would have a larger marginal utility for low video rates in comparison to higher ones. With users viewing video on a range of devices with varying capabilities, the video rate utility would correspondingly vary across heterogeneous devices. One option to account for this heterogeneity is to tune ρ_u by the operator to differentiate the quality of service provided to

users with different priorities. Given the exponential utility function is upper-bounded by 1, ρ_u can be estimated as

$$\rho_u = -\frac{r_q}{\overline{r_u}} \log(\epsilon), \tag{2}$$

where $\overline{r_u}$ represents the maximum rate assigned to user u due to device capability or assigned user priority. With ϵ being a small fraction, this design implies that $\Upsilon_u(\overline{r_u}) = 1 - \epsilon$ and that any increase in the streaming rate beyond $\overline{r_u}$ for user u would lead to an insignificant change in the user utility.

SAP seeks to capture the stall impairment by estimating the probability of buffer depletion at the client. User u would typically stall if the download time of the requested segment, denoted as d_u, is larger than the segment deadline D_u. D_u can be estimated as the time remaining until segment playout. Since segments are typically downloaded sequentially, the next segment deadline equals the buffer-level at the time of sending the segment request. The download time d_u depends on the downloaded segment size and the delivery rate of this segment to the client. The downloaded segment quality is independently selected by the client according to its adaptation policy based on the application state information. Assuming that the wireless access link is the bottleneck, the delivery rate of the segment would mainly depend on the resource scheduling at the BS, which of course only considers network state information when it is allocating resources to its users.

In essence, SAP indirectly controls the BS scheduling, by managing the flow of packets toward the BS, considering both network and application states. This is performed at an intermediate timescale between the small BS scheduling period and the much longer quality adaptation time scale. Hence, when the SAP pacing manager is executed, users can be classified into one of two states

• new-request state in which the user has recently finished downloading a segment and is requesting a new one, or

• mid-request state in which the user is still downloading a previously requested segment.

With the first class, SAP determines the stall probability of user u, denoted as π_u, as

$$\pi_u \quad = \quad Prob(d_u > D_u) = Prob\left(\frac{S_{\vartheta_u}}{x_u} > D_u\right),$$

where x_u represents the rate at which user u's packets would be delivered to the BS and S_{ϑ_u} represents a segment size random variable conditioned on the selected quality level $\vartheta_u \in \{1..Q\}$ corresponding to the selected x_u. It is worth noting that for the new-request state, x_u represents the quality rate of the new segment that should ideally be requested by the client from SAP's point of view. This quality identifier ϑ_u is maintained by SAP as part of the flow state until it is updated with the next segment request. Hence, if the system state made SAP to select a high quality rate at the beginning of the segment, this decision would be supported until the next segment request, even when network state changes.

If the user is in the mid-request state, the SAP pacing manager calculates the stall probability using the conditional residual segment size distribution, denoted as $F_{S_{\vartheta_u}}(.)$, and is expressed as

$$\pi_u = \frac{1 - F_{S_{\vartheta_u}}(b_u + D_u x_u)}{1 - F_{S_{S_{\vartheta_u}}}(b_u)}, \tag{3}$$

where b_u represents the total transmitted bytes for the currently downloaded segment of user u. Note that segment download deadline D_u is reduced by the time elapsed since the last execution of the SAP pacing optimization. Clearly, there is a need to monitor both downlink and uplink. The downlink is monitored to determine the number of bytes transmitted per segment for each active user. The uplink is monitored to identify new segment requests.

3.2.3 SAP Optimization Program. SAP maximizes the total quality utility minus the total stall penalty for all users. This design is inspired by the desire to provide an improvement in the QoE performance of video delivery on the cellular network, and particularly to seek consistent QoE for users, even those with diverse wireless channel conditions. Our SAP pacing manager optimization program is expressed as

$$\max_{x_i} \sum_{u=1}^{U} \Upsilon_u(x_u) - \beta \pi_u$$

such that

$$\sum_{u} \quad x_u/\gamma_u < \zeta \quad C \tag{4}$$

$$x_u \qquad \in \qquad \{\hat{r}_1, ..., \hat{r}_Q\} \tag{5}$$

where β represents non-negative weights for the switching and stall penalties, respectively. $\hat{r}_i$ represents the scaled version of the video encoding rate, and ζ represents a scaling factor to avoid resource underutilization at the BS. We scaled the encoding rates r_i to $\hat{r}_i$ to compensate for the overhead of lower layers such that the application goodput would match the target encoding rate. Additionally, we employed the scaling factor ζ to avoid resource underutilization, resulting from performing SAP rate control using the reported wide-band spectral efficiency while the scheduling performed at base station is based on individual resource unit spectral efficiency.

The integration of quality and temporal components in the SAP optimization assist streaming clients to avoid stalls, irrespective of the cause of stalls. Stalls may happen for different reasons including significant changes in the user channel condition, sudden changes in network load (e.g., arrival of new users), and/or large increase in the video's bandwidth demand due to the inherent variable bitrate of compressed video. SAP accommodates many of these scenarios by dynamically re-adjusting the resource allocation and thus controlling the packet delivery process to the BS scheduler. By throttling a specific user, the BS is implicitly forced to serve the traffic of other users.

The operation of the SAP pace optimizer depends on the availability of a set of conditional distributions for segment sizes using different encoding rates for a given segment duration. These can be easily determined by fitting the segment size data to a suitable distribution. A set of such distributions are prepared in advance by the network operator for a wide

Table 1: Weibull distribution parameters for segment sizes

quality	1	2	3	4	5	6	7	8	9	10
Shape	2.65	2.7	2.72	2.84	2.9	3.18	3.07	3.21	3.15	3.20
Scale	132	210	313	419	586	973	1309	1666	2140	2388

range of content, with segment size data made available using one of several techniques. One option is to fetch byte-range MPD files if they are available; these provide segment sizes for each encoding rate. Or segment sizes can be obtained by iteratively downloading segments at several fixed qualities from the content provider and recording the segment sizes. Alternatively, if there was cooperation with the content provider, segment size data (or indeed distributions) could be provided to the network operator directly. In this work, we obtain segment size data from the full iVID dataset [24] and fit this data to a Weibull distribution using fitdistplus package in R.

Table 1 shows the scale and shape parameters of the fitted Weibull distributions for different quality levels of all movies in the publicly available iVID dataset. The segment duration for these is 4 seconds.

The SAP pacing optimization program can be classified as a non-linear discrete optimization problem. Solving such problems is usually time consuming and not necessarily feasible in real-time. However, we ensure real-time operation by taking advantage of the problem structure. Since all nonlinear terms are functions of a single optimization variable, the program can be formulated as a *separable programming* problem that can be solved at a speed similar to linear programs [9, 12].

3.3 Collaborative vs. Non-collaborative SAP

The previous section detailed the design of SAP and how it makes use of network and application states. SAP can operate in one of two modes: collaborative or non-collaborative. In collaborative SAP, the application state logger acts as a DANE interface that receives status messages from the client. This provides SAP with encoding information (rates and segment duration) and the client's current buffer-level. Note that the segment duration is used only for identifying the conditional segment size distribution to apply. Encoding parameters would be provided at the beginning of the session, while the buffer-level would be reported on a regular basis during the session. In the non-collaborative case, we assume that the network operator would implement additional in-network functions to identify the required information for SAP. We now explore the non-collaborative case and consider two possibilities:

• ***Non-encrypted client-CDN communication.*** In this case, the operator would be able to extract encoding parameters, such as encoding rates and segment duration, from the DASH description (MPD) file. To estimate the buffer-level, the SAP application state logger would need to rely on deep packet inspection. Generally, the uplink mainly carries HTTP GET requests for video segments and TCP ACK packets. By tracking HTTP requests, SAP can estimate a conservative value for the client application buffer as detailed below.

• ***Encrypted client-CDN communication.*** In this case, SAP has to rely on *a priori* information about video traffic. General guidelines for encoding rates used by different video content providers are publicly available. Hence, SAP may rely on these public rates as alternatives for r_i. Note that the SAP pacing optimizer may use arbitrary discrete rates for r_i with inter-rate gaps following the public guidelines. A typical rate ratio between consecutive rates is 1.5. But, we believe that the minimum rate is the most critical factor. Choosing the right value for the minimum rate is important, especially for users with poor link conditions: choosing too small a value may over-throttle the flow to the user, resulting in excessively poor QoE; on the other hand, choosing too large a value may result in inefficient utilization of the available resources. For our work in this encrypted scenario, we assume a technique to estimate the segment duration by observing the segment size of the first few segments. This technique may assume a small segment duration (e.g., 4 sec) until a more accurate estimate of the segment duration is calculated. Similar to the non-encrypted case, this segment duration would be used to estimate the buffer-level as presented below.

3.3.1 SAP Buffer-level Estimator. The buffer-level estimation is needed in the non-collaborative scenario to identify the delivery deadline of segments. The buffer-level, denoted as b, can be estimated as the difference between the received and played video durations, denoted as D_r and D_p respectively. Hence, the buffer-level can be calculated as

$$b = D_r - D_p. \tag{6}$$

By knowing the segment duration and the number of received segments, the received media duration can be directly estimated as their product. The number of received segments can be determined by monitoring the HTTP GET requests on the user uplink, denoted as g_u for user u. The received video duration may be estimated as $(g_u - 1)\tau$ seconds. It is important to note that this approach would efficiently work in both non-encrypted and encrypted cases, provided that the client requests segments in a sequential manner. However, if the client opts to abandon segments while they are being downloaded, for example to request lower quality segments to avoid stalls, then g_u should correspond to sequential segment requests. Note that in such mechanisms today the client is required to close the TCP connection to trigger abandonment of the requested segment and establish a new connection to request the new segment. Hence, this exchange can be captured to maintain an accurate estimate of g_u.

The playout video duration can be estimated as the difference between the current wall clock, denoted as t_c, and the time at which playout starts, denoted as t_p. The initial buffering depends on the client implementation, but an initial buffering duration between 6 and 10 seconds is common and considered an acceptable startup delay for users. Our estimator assumes that the client would always start after

downloading the first segment. The validity of this assumption is based on the observation that larger segment durations are used for cellular clients. Hence, the estimated buffer-level would be close to the actual buffer-level. However, it would be a conservative estimate for the buffer-level if the client actually has a larger initial buffer. Additionally, this estimate would deviate from the actual buffer if the user intervenes with the session, e.g., pause the playout. Note that in this case, the client would continue to download the media until the buffer saturates and then would stop requesting segments until the playout is resumed. In this case, SAP is agnostic to user action and would consider all clients as actively playing out their entire buffer.

Since we depend on the real-clock to capture the play-out duration of the stream, the duration of an interruption should be used to rectify the buffer-level estimate. In our estimator, stalls are captured by negative buffer estimates. To rectify the buffer-level estimate, we shift our playout reference time to the instant at which the client resumes playout after a stall. We assume that the client would continue the playout once it receives a segment. Additionally, we reset the number of received segments to 1 (i.e., $g_u = 2$). Hence, we can consider t_p as the time at which playout starts or resume after a stall. Similarly, g_u can be considered the number of stall-free received segments since the client started or since the last stall.

4 PERFORMANCE EVALUATION

4.1 Evaluation Setup

The SAP evaluation testbed is designed based on an empirical methodology that requires the use of realistic video content and player software, operating over a simulated cellular network, thus promoting realism and enabling controlled repeatability of experiments across a range of configurations. Fig. 2 shows our evaluation testbed, which is based on iVID D-LiTE testbed [23]. In the following subsections, we present the key elements of our testbed and we provide full implementation details in the appendix.

4.1.1 Streaming Clients. In our testbed, DASH clients are running over a version of Ubuntu installed in Raspberry Pi 2 and standard net-books. The clients use GPAC 0.5.2-DEV-rev985[4] that we extend with well-known/recent adaptation algorithms such as BBA2 [14], FESTIVE [15], conventional (CONV) [15, 16], and ARBITER [29]. BBA2 represents the class of buffer-based algorithms in which the quality selection mainly depends on the current buffer-level. FESTIVE represents the class of conservative rate-based algorithms that are designed to operate well in scenarios with shared bottlenecks. CONV represents another rate-based strategy that streams video while trying to maintain the quality at a stable level as much as possible. ARBITER represents the class of hybrid algorithms that integrate both application and network state in the quality selection decision. The parameters of different streaming algorithms are set to the default values reported in the papers cited. In all the evaluation scenarios, a client is configured to perform 8 seconds (two segments) of initial buffering and 4 seconds (one segment) of rebuffering after any stall [26].

4.1.2 LTE Network Setup. The LTE network is implemented using the LTE-EPC network simulator (LENA) module[5] in ns3[6]. In our setup, external nodes are connected to simulated nodes using the ns-3 TAP mechanism, which uses a special net device called a TapBridge. In order to connect LTE UEs to the external streaming clients, a second carrier sense multiple access (CSMA) net-device (12.0.0.x network) is added to the UE, as LTE-net-devices are not compatible with the ns3-TAP bridge. On the other side of the LTE network, the LTE packet gateway (PGW) is connected to the network attached storage (NAS) server through a master node whose functionality is explained below. We have modified the code of both the LENA scheduler and routing modules in ns3 as described below.

For proper routing of downlink packets, we added static routes for the 12.0.0.x network at the remote node. Additionally, we changed the implementation of ipv4-list-rout-ing.cc to mangle the destination address of downlink packets before entering and after exiting LENA devices to allow the packets to travel through the tunnel between LTE PGW and UEs. The destination address is changed from 12.0.0.x to 7.0.0.x before being forwarded to the PGW. Additionally, the destination addresses of downlink video packets received at an LTE netdevice (7.0.0.x network) are changed back to the actual client 12.0.0.x address and are then forwarded to the ns3 12.0.0.x TapBridge. For uplink traffic, the packets follow the default gateway towards the PGW and are then forwarded based on static routes that are installed at the LTE PGW and remote host node for the NAS network (9.0.0.x) to be forwarded to the appropriate device. Each physical client node (Raspberry PI or netbook) is configured with a default gateway which is the corresponding CSMA netdevice of the connected LTE UE. The packets then proceed to the LTE PGW as the default gateway for LTE UEs.

In LENA, we consider the log distance path-loss channel model [8] for the link between eNodeB and UEs. This pathloss is overloaded with fading traces generated using a tool provided with LENA. The default LENA configuration parameters are used for both eNodeB and pathloss model. All our evaluations are conducted with the proportional fairness (PF) scheduler at the eNodeB. We have modified the PF scheduler implementation only so that user channel quality can be reported to the traffic manager being evaluated.

4.1.3 Master Controller. The master controller node is responsible for orchestrating the evaluation and performing traffic management functions during video sessions. More specifically, this node configures the network and GPAC clients before an evaluation run starts. Note that network configurations include parameters such as eNodeB bandwidth,

[4] https://gpac.wp.mines-telecom.fr/

[5] LENA module. https://goo.gl/6D1Wfq. Last accessed: Dec 8, 2016.

[6] https://www.nsnam.org/

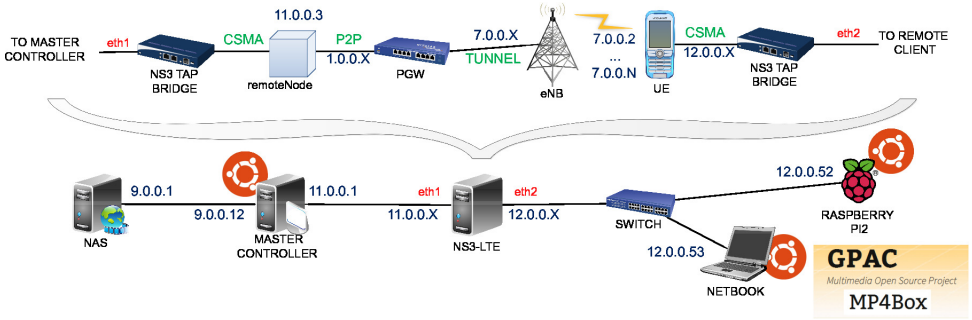


Figure 2: Evaluation Testbed

eNodeB scheduler, and fading model, while the client configuration parameters include the adaptation algorithm and streamed video specifications. At the end of the run, it also collects performance logs from the clients for post processing.

During the evaluation, the master controller performs traffic management functions. In our evaluation, we compare the performance of three different traffic controllers including no traffic control (noTC), AVIS [4], and SAP. AVIS focuses on fair user allocation and reduction of quality switches, but overlooks stalls in its design. Both AVIS and SAP optimization programs are implemented using the Lindo solver[7]. The AVIS implementation is based on the continuous version presented in [4] and the SAP implementation is based on the separable programming model. In our testbed, the AVIS parameters are set to the default values presented in [4]. For SAP, we set the stall weight β to 100 and the default utility saturation parameter ρ_u of all clients to 3Mbps. 3Mbps is the maximum rate for HD resolution in the iVID video dataset.

Both SAP and AVIS are periodically executed every 250 ms and the outcome of the optimization program is used to throttle individual user queues using the traffic control command in Linux. The master node maintains a communication channel with the eNodeB to obtain user channel quality information (CQI) required by both AVIS and SAP. Similarly, it maintains communication channels with streaming clients to obtain relevant application information for the collaborative scenarios. In the non-collaborative scenario, the buffer-level is estimated as presented in Section 3.3.1.

4.1.4 Evaluation Scenarios. In our evaluation, we consider a group of video users sharing a highly loaded LTE eNodeB to capture the impact of streaming in a limited-resource environment. Our performance evaluation shows that different traffic management solutions have similar performance in lightly loaded scenarios. Hence, we focus more on the more relevant highly loaded systems. The eNodeB bandwidth is 1.4 MHz and has a transmission power of 30.2dBm. This eNodeB has 6 resource units split into 6 allocatable resource block groups (RBG) in the downlink. Each RBG can support a PHY rate between 16 Kbps and 712 Kbps depending on the user reported channel quality indicator value. In our evaluation, we consider two different user topologies in highly loaded cells as detailed below.

[7]http://www.lindo.com/

We consider 6 DASH clients streaming 5-minute, 4-second segment, videos over a single LTE eNodeB. In every session, the clients are introduced to the network separated by a 1 second time gap. Each user streams a different five-minute video from the iVID dataset [24][8], whose videos are encoded with 4-second segments at the following rates {235, 375, 560, 750, 1050, 1750, 2350, 3000, 3850, 4300} Kbps. Note that the base station is only used by the video users. In a more general setup, other traffic may share the eNodeB and video users would be allocated a slice of the total base station resources.

4.2 Performance metrics

Our performance metrics include the average received data rate per session (r_{av}), the average number of stalls per session (n_{st}), the average stall duration per session (t_{st}), the average number of switches per session (n_{sw}), the average switching level (l_{sw}), and a combined QoE metric (x_q) [20], which was originally developed in [17] and [27]. The QoE metric is expressed as [20]

$$x_q = max(0, 0.17 + 5.67\frac{q_{av}}{q_Q} - 6.72 * \frac{q_{std}}{q_Q} - 4.95\varphi),$$

where q_{av} and q_{std} represents the average and standard deviation of the received quality level, and φ represents the stall penalty and is expressed as

$$\varphi = 0.875 * max(0, 1 + ln(f_{st})/6) + 0.008333 * min(t_{st}, 15),$$

where f_{st} represents the frequency of stalls and t_{st} is the average stall duration per session. The results shown represents the average of each metric obtained across 15 runs.

4.3 Performance Results

4.3.1 Scenario 1: Collaborative Clients with Diverse Link Conditions. In this scenario, we consider equally separated users with the nearest and farthest users being at 25m and 375m, respectively, from the base station. Fig. 3 illustrates the distribution of the achievable rate per resource unit for each user. Note that these rates are determined by mapping the reported CQI value to the eNodeB as defined in the LTE standards. The figure shows that the closest client to the eNodeB can achieve the highest PHY rate per RU (712 Kbps) almost all the time, while the farthest client's achievable rate is less than 200Kbps for 60% of the time. Fig. 4 shows

[8]iVID Dataset. https://goo.gl/BP6LWR. Last accessed: Dec 8, 2016.

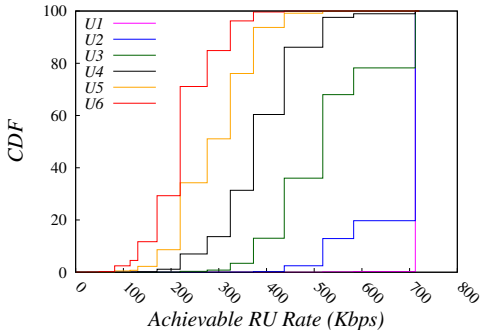
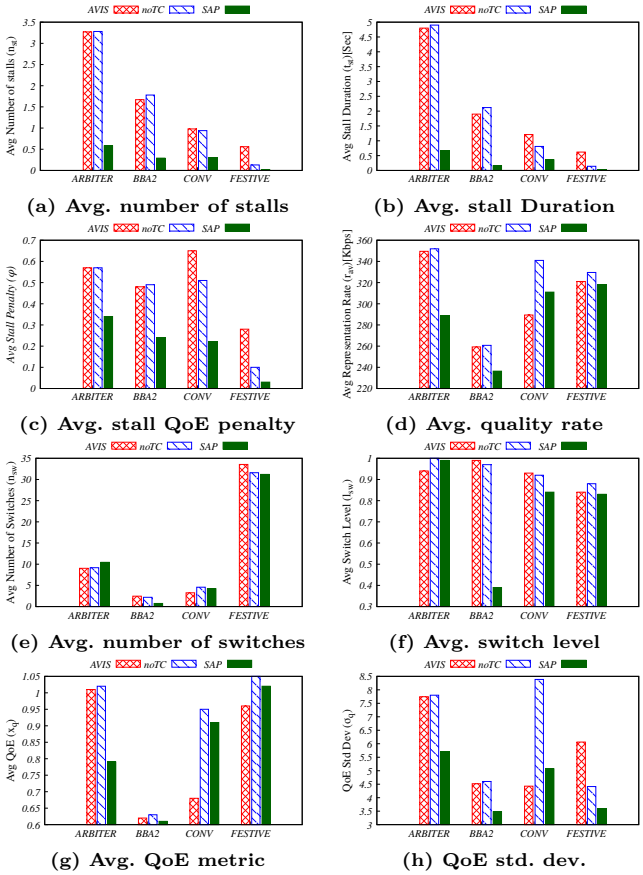


Figure 3: Achievable rate per resource unit (RU) per user in scenario 1

the average of each of the performance metrics for different streaming algorithms.

SAP significantly improves the stall performance for all algorithms in comparison to both noTC and AVIS. The existing diversity in the channel condition enabled SAP to achieve a significant reduction in both the number of stalls, n_{st}, and stall duration, t_{st}. In conjunction with BBA2, SAP reduces both n_{st} and t_{st} by 84% and 94% in comparison to noTC. SAP also reduces both n_{st} and t_{st} of FESTIVE by 95%. Additionally, we note a significant reduction of 83% and 86% for n_{st} and t_{st} when SAP is used with ARBITER. Similar performance gains are achieved by SAP in comparison to AVIS. In fact, AVIS shows stall performance that is similar to or worse than noTC. The improvement in stall performance by SAP thus reduces the stall QoE penalty φ by 50%, 89%, 66%, and 40% for BBA2, FESTIVE, CONV, and ARBITER, respectively. In this scenario with users having diverse link conditions, the user farthest from the eNodeB encounters most of these stalls.

The improved stall performance for SAP is due to its resource management strategy, that not only ensures a minimum rate for every client, but also dynamically paces individual stream packets to protect clients from stalling. Fig. 5 illustrates the resource shares allocated by SAP and AVIS for individual streams. Clearly, Fig. 5a and Fig. 5b illustrates that in SAP more resources are allocated to protect the farthest client who is more exposed to stalls. It is also interesting to observe that with BBA2, SAP is continuously supporting this user by allocating more resources over the entire session duration, in comparison to only the initial part with FESTIVE. Note that BBA2 employs a large buffer (240 sec) and its clients are continuously competing for the network resources. On the contrary, FESTIVE employs a much smaller buffer (30 Sec). Hence, FESTIVE clients with good channel conditions are able to fill their buffer and consequently delay their segment requests. Thus, clients with poorer channel conditions are offered more transmission opportunities by the eNodeB and SAP stall avoidance would kick in at a much lower frequency in comparison to BBA2. On the other hand, Fig. 5c shows that while AVIS tends to avoid frequent rate changes it does not provide the same level of stall protection to video clients.



(a) Avg. number of stalls (b) Avg. stall Duration

(c) Avg. stall QoE penalty (d) Avg. quality rate

(e) Avg. number of switches (f) Avg. switch level

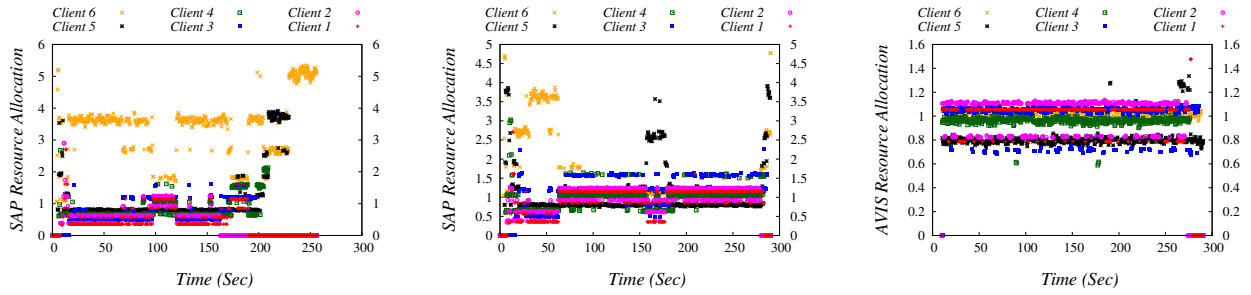
(g) Avg. QoE metric (h) QoE std. dev.

Figure 4: Performance metrics for scenario 1

Fig. 4d shows that introducing a network traffic controller, e.g., AVIS or SAP, results in reducing the average representation rate r_{av}. This reduction is expected, as pacing the traffic would slow the delivery rate of packets, leading to lower rate estimates for rate-based strategies or lower buffer-levels in buffer-based strategies. Fig. 4e and Fig. 4f show that SAP significantly reduced the average number of switches n_{sw} and average switching level l_{sw} of BBA2 by 64% and 60%, respectively. Excluding BBA2, the switching performance of the remaining algorithms is not significantly affected by the traffic pacing/control functions.

Fig. 4g shows that the QoE metric x_q drops when SAP is used. This reduction is due to SAP reshuffling resources to help suffering clients - which are those close to the cell edge. Such reshuffling reduces x_q for users close to the eNodeB and increases x_q for distant users.We emphasize that SAP provides improved fairness, by reducing the variance in QoE as shown in Fig. 4h. The latter figure shows that SAP reduces σ_q by 40%, 40%, 24%, and 26% in comparison to noTC for BBA2, FESTIVE, CONV, and ARBITER, respectively.

4.3.2 Scenario 2: Collaborative Clients with Similar Network Conditions. In our second scenario, we consider users with a mobile vehicular fading channel at a distance of 300m from the base station. These clients have an average achievable



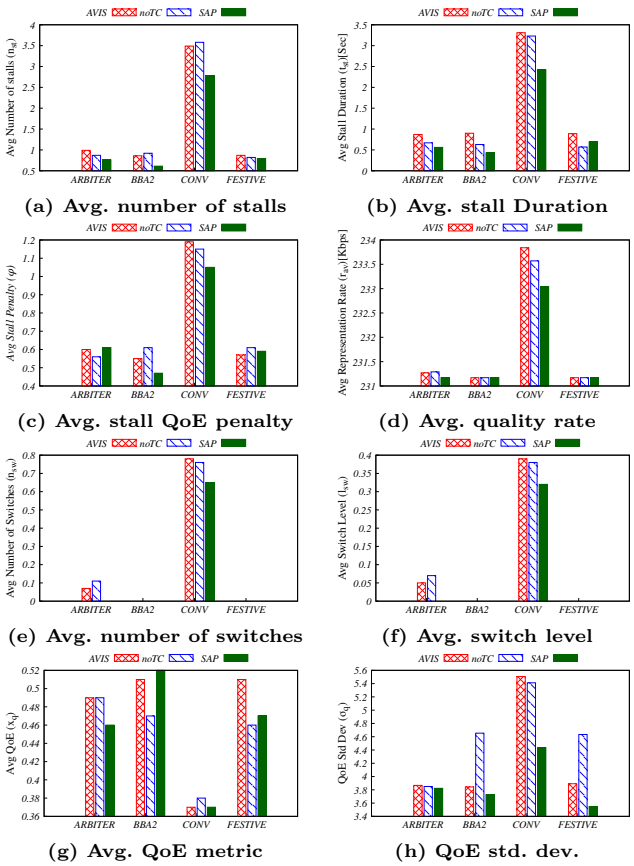
(a) BBA2-SAP (b) FESTIVE-SAP (c) FESTIVE-AVIS

Figure 5: Comparing SAP and AVIS resource allocation with different algorithms

wireless RU PHY rate of 301Kbps that over time varies between 16Kbps and 712Kbps. Fig. 6 plots our key performance metrics for different adaptation strategies and different traffic management strategies. Fig. 6d-6f show that the traffic controller has limited impact on the quality selection for both FESTIVE and BBA2, although SAP does reduce the representation rate of the streamed video for reasons explained earlier. But, SAP does reduce the number of switches, and switching level for the ARBITER and CONV algorithms.

Fig. 6a shows that SAP manages to noticeably reduce the number of stalls encountered for all adaptation strategies in comparison to both AVIS and noTC. This reduction reaches up to 34% and 29% in comparison to noTC and AVIS respectively. Fig. 6b shows that SAP also reduces the average stall duration per session for most of the combinations of adaptation strategies and traffic controllers. This improved stall performance leads to a reduction in the stall penalty up to 23%, 9%, and 3.5% for BBA2, CONV, and FESTIVE, respectively. For ARBITER, we found that SAP's reduced stalls are in fact spread across more sessions leading to this increase in the stall penalty of ARBITER in comparison to noTC. Note that such a distribution of stalls across users leads to a reduced QoE metric as shown in Fig. 6g but reduces the variance in the QoE observed by the users sharing the same resources as shown in Fig. 6h. Note that SAP helps both FESTIVE and BBA2 clients to boost the QoE and reduce the QoE variance among competing clients when compared to noTC.

Fig. 7a plots a selected subset of performance metrics for individual BBA2 clients with both noTC and SAP. Note that BBA2 clients persist at the lowest quality. Because of this, showing quality rate and switching metrics for this would not be useful. The figure illustrates that SAP achieve the performance gains by significantly reducing the stalls encountered by the second client. We noticed that the second user stalls more than the other users because its video rate increases for a series of subsequent segments, leading to buffer depletion and a series of stalls. To illustrate, Fig. 7b and Fig. 7c plot the dynamics of the video client for user 2, with the noTC and SAP traffic controllers, respectively. Clearly, we see that the video client sticks to the first quality representation for the entire video. However, the video demand increases mid-session, and for a series of subsequent segments, leading to buffer depletion as these segments take longer time to

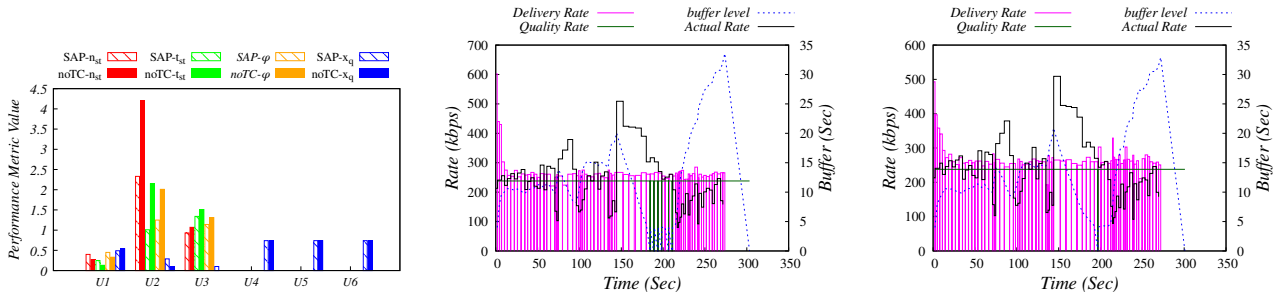


(a) Avg. number of stalls (b) Avg. stall Duration
(c) Avg. stall QoE penalty (d) Avg. quality rate
(e) Avg. number of switches (f) Avg. switch level
(g) Avg. QoE metric (h) QoE std. dev.

Figure 6: Performance metrics for scenario 2

download (note the wide bars in the delivery rate curve). Such a situation is handled well by SAP because it allocates more resources to this client. Thus, SAP helps to reduce the number of stalls from 6 with noTC to just one stall for the session shown. Note that the delivery rate and quality curves (b-c) are not time-aligned as delivered segments are stored in the buffer before being played out. Additionally, the quality rate curve captures the stalls by dropping to zero for the duration of the stall, if any.

The integration of application and network states enables SAP to assist users as needed when network conditions or video demands change. With static users, edge-users are disadvantaged. SAP assists edge-users by optimally throttling



(a) Individual BBA2 client performance with noTC and SAP (b) BBA2 client dynamics over noTC (c) BBA2 client dynamics over SAP

Figure 7: Detailed client performance and session dynamics example in scenario 2.

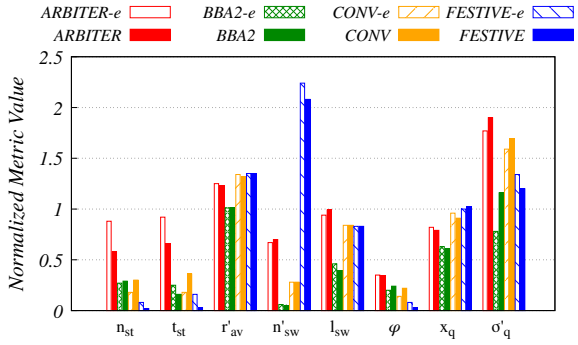


Figure 8: Normalized streaming performance metrics for collaborative and non-collaborative scenarios. -e denotes the non-collaborative version. $r'_{av} = r_{av}/235$, $\sigma'_q = \sigma_q/3$, and $n'_{sw} = n_{sw}/15$.

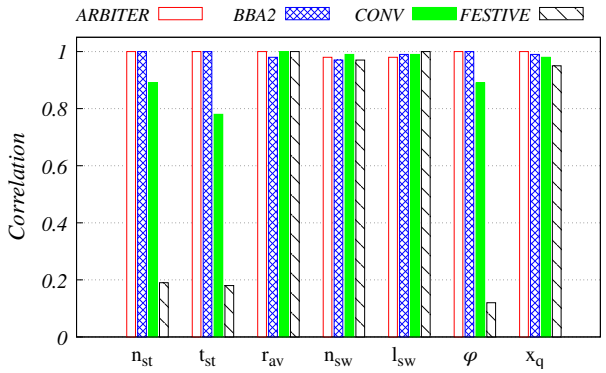


Figure 9: The correlation between the performance metrics of collaborative and non-collaborative scenarios

the traffic of other users as shown in Scenario 1. SAP can also support users as video demand increases as shown in Scenario 2. We believe also that SAP can assist users that experience large scale variation in link conditions due to user mobility. Such changes can result in a noticeable drop in the user data rate that would lead to a longer segment download time. Hence, SAP would proactively help users with a degrading channel, especially those users which SAP estimates as having a low buffer-level. This assistance can avoid stalls and large quality changes. Hence, we believe that SAP would benefit both static and mobile scenarios.

4.3.3 Collaborative vs. Non-collaborative SAP. In this section, we investigate the performance of SAP in non-collaborative mode in which the buffer-level is estimated by SAP in comparison to being collaboratively provided by the client. In the non-collaborative mode, we consider that encoding information is available to SAP, but the instantaneous buffer-level has to be estimated by SAP, using the estimation technique presented in Section 3.3.1. Fig. 8 shows the streaming performance metrics averaged over all sessions. Note that we normalized some of the metrics to fit all of them in one figure, and the normalization factors are indicated in the caption. This figure shows that both collaborative and non-collaborative SAP achieve similar performance metrics for all the adaptation algorithms. Additionally, Fig. 9 plots the correlation between individual client performance metrics for both collaborative and non-collaborative scenarios. The figure shows a high correlation for all metrics for all adaptation strategies except for the FESTIVE stall metric. The low correlation for the FESTIVE stalls is due to a minor variation in the stall distribution across the clients. Only the farthest client encounter stalls in the collaborative scenario while in the non-collaborative scenario the second farthest client also has a couple of stalls. Hence, we conclude that integrating the buffer estimation technique with SAP can be sufficient, thus avoiding any frequent updates from the streaming end client. Consequently, the network does not need to have a direct interface with client software or be aware of the local state of the individual client adaptation.

5 CONCLUSIONS

The ability of DASH video streaming clients to operate effectively over cellular networks is poor, due to the inherent variability of both video datarates and the wireless channel quality. In highly loaded systems, these characteristics may lead to streaming issues, such as stalls. Unlike the situation when using wired links, it must be recognized that achieving equivalent QoE across all clients is not a reasonable objective. Fairness is not absolute, but is rather a function of the channel conditions available to each client. Thus, steps must be taken to manage wireless resources to ensure that clients with poorer channel conditions do not suffer unnecessarily, even as clients with good channel conditions seek to maximize their own QoE. This can be achieved by the network redistributing channel resources so as to reduce extremes of QoE

across clients. SAP seeks to re-allocate wireless resources in favor of a client that becomes less likely to experience a stall, without significantly degrading the QoE of other users. In this paper, we present SAP as a traffic management solution to improve the streaming performance of a group of DASH video users competing for a base station's resources. SAP integrates both application and network state to optimally pace individual stream delivery. Our extensive experiments show that SAP significantly reduces video stalls encountered by different users in comparison to the state of the art solutions. Additionally, SAP reduces the QoE variability across multiple clients, leading to a more consistent user experience.

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6 APPENDIX: TESTBED IMPLEMENTATION DETAILS

In this Appendix, we present the steps required to replicate our testbed setup. Our testbed is based on iVID D-LiTE testbed, which is developed as part of the iVID project at University College Cork, Ireland. Our testbed is designed based on an empirical methodology that uses real end-points operating over a simulated network, LTE network in our case. This model promotes realism and enables controlled repeatability of experiments across a range of configurations. We extend D-LiTE to test the performance of different streaming algorithms while performing traffic management functions, which may implicitly include interfacing the controller with other network nodes to obtain application and network state information. Fig. 2 shows our evaluation testbed. In the following, we present implementation details for the key parts of our evaluation testbed. Additionally, we provide an online step-by-step guide[9] for building different testbed elements.

6.1 Testbed Setup

6.1.1 Video Content and Server Setup. Our video content is based on the iVID dataset [24]. DASH uses standard HTTP servers that host the video content in its content directory. In order to setup the server, one should download the dataset content into the content folder of the selected HTTP server. For example, the Apache default content folder in Linux is /var/www.

6.1.2 Streaming Client Setup. Our streaming clients are based on GPAC, which is installed in a typical UBUNTU desktop. In the SAP testbed webpage, we provide a script that can be used to setup GPAC and its dependencies for both H.264 and H.265 on an UBUNTU desktop. Note that we use static IP addresses for both physical and virtual interfaces.

The released GPAC code implements different streaming algorithms including BBA2, FESTIVE, ARBITER, and CONV algorithms. For experimentation purposes, we also modified GPAC behavior to suit different algorithms. Specifically, we changed the buffering and ON/OFF behavior of GPAC.

GPAC implements a two-level buffering hierarchy including cache and playout buffer. The cache directly receives from the socket while the playout buffer loads media segments from the cache as needed by the decoder. The default behavior in GPAC is to fill the cache buffer before decoding the media. Our modifications enable setting the initial playout delay, rebuffering delay (in segments), and the cache and playout buffer sizes in milliseconds.

Further, we modified the default ON/OFF behavior. By default, GPAC stops segment download when the cache buffer is full and resumes when it drops to half its size. This behavior is overridden as request scheduling is now a design component of different streaming algorithms. For example, FESTIVE considers a random delay before sending the next request. Hence, an algorithm-dependent request delay is implemented by introducing a request_delay variable, determining the next GET request delay in milliseconds.

Our main GPAC modifications extend over five files including

- *GPAC.cfg* is extended to include algorithm-specific and general parameters.
- *export.cpp* defines additional template links for these parameters.
- *mpd_in.c* is changed to determine the algorithm according to the parsed configuration.
- *dash.h* includes headers *for functions* used for parsing the configuration.
- *dash_client.c* includes the implementation of configuration parsing functions and the adaptation logic of different algorithms.

GPAC calls the adaptation logic from a function called *dash_ adaptation_selection*. We also implemented helper functions, such as calculating the first and second order statistics of throughput samples and getting buffered media duration, that are shared by different algorithms. In order to implement a new adaptation algorithm, one should introduce modification to all these files.

Note that, some algorithms, such as BBA and ARBITER, require segment size information in their decision. Byte-range MPD files include this information while URL based MPD files do not. In our implementation, such data is made available to the client by maintaining a local copy of segment size information provided by the third iVID Dataset [24]. Note that if another dataset is used, one should prepare segment size information files in order to use algorithms such as BBA and ARBITER.

6.1.3 Network Simulator Setup. The network simulator node is used to emulate LTE while interacting with real end nodes, i.e., real client and servers. Real and simulated nodes are connected using a network consisting of virtual and real network interfaces. As shown in Fig. 2, every real client is connected to a real switch that is connected also to one of two Ethernet interfaces of the simulator node. Every simulated user equipment (UE) has two netDevices including one connected to the LTE network and a CSMA netDevice connected to a Linux TAP device. UE TAP devices are connected to a Linux bridge device that is also connected to the Ethernet interface connected with the real clients. Hence, the simulated nodes and corresponding real clients are connected using an Ethernet network consisting of a virtual bridge device and a real Ethernet switch. In the SAP testbed webpage, we provide shell scripts that can be used to setup (bridge_network_build.sh) and tear down (bridge_network_teardown.sh) the required virtual devices. We also provide our LTE simulation code to setup the simulated network and connect UEs to corresponding TAP devices (lena-simple-epc-tap-nas.cc). Note that one should create internal TAP and bridge devices before running the simulator. Proper routing and packet mangling rules for uplink and downlink should be set up as described in Section

[9] https://www.ucc.ie/en/misl/research/software/sap-testbed/

4.1.2. We also provide a modified ipv4 routing module file (ipv4-list-routing.cc) for ns3 in the SAP testbed webpage.

6.1.4 Master controller Setup. The master controller node performs three main functions. It orchestrates the experiments, interfaces with other nodes for obtaining relevant information, and performs traffic management functions.

Experiment Orchestration . The master controller orchestrates the experiment by initializing other testbed elements (streaming clients and network), terminating the experiment, and collecting performance log files. Before the actual experiment starts, the master controller configures the streaming client with the adaptation algorithm, buffer size, and other relevant parameters. These parameters are passed to a shell script that is distributed in all clients. This script takes responsibility of preparing the streaming client configuration file before starting the video client. The configuration file includes information about the selected adaptation logic and relevant algorithm parameters. Additionally, the master controller configures the network node with relevant mobility and channel configurations. Note that the simulator code (lena-simple-epc-tap-nas.cc) is developed to parse these configuration and set up the network accordingly. On configuring all nodes, the master controller uses distributed shell utility (*dsh*) in Ubuntu.

Controller Interface. The master controller maintains a control channel with the simulator and streaming nodes to obtain relevant information for traffic management functions. For example, the traffic management may need information about the link quality of individual clients and act accordingly. Additionally, streaming clients may exchange application state information with the network controller over a different control link. In our implementation, the communication between the master controller and other nodes takes place over the distributed shell used for orchestrating the experiment. Each video client and the simulator node log the required corresponding information to the shell. Exchanged information over such channels are parsed for specific information, such as channel quality and buffer-level at the client.

Traffic Management. In our testbed, we implement per flow traffic rate control for both AVIS and SAP. Both solutions employ an optimization framework that identifies the target rate for individual sessions. The rates are enforced using typical token bucket filter (*TBF*) in the traffic control (*tc*) command in Linux. The traffic from all sessions is also limited to the total of the estimated individual flow rates. Note that both AVIS and SAP avoid resource waste by allowing the usage of the unused residual bandwidth up to a limit identified as the average rate of the selected quality rate and next quality rate. In the iVID D-LiTE web-page, we provide the master controller shell scripts. Note that the code is distributed over multiple directories and files as described in the README file.

To implement a new traffic management function, one should introduce a new branch for the if statement handling different traffic management functions.

6.1.5 Testing your setup. We would like to stress that ping packets do not travel through the simulated LTE network. Hence, one should test the connectivity using other applications, such as web browsing or iperf, running at the physical devices.